# Crypto Part 3: Certificates & TLS CMSC 23200, Winter 2024, Lecture 6

Grant Ho & Blase Ur

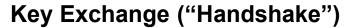
University of Chicago (Slides adapted from David Cash)

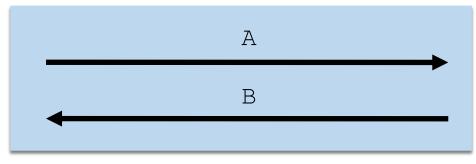
## Outline: Crypto Part 3

- Recap: Secure communication channels
- Authenticating endpoints: Certificates (Certs)
- Issuing Certs and Certificate Infrastructure (PKI)
- Attacks, Countermeasures
- Real World Secure Channels: SSL / TLS

### Template For Secure Channels (TLS, SSH, IPSec, ...)







uchicago.edu



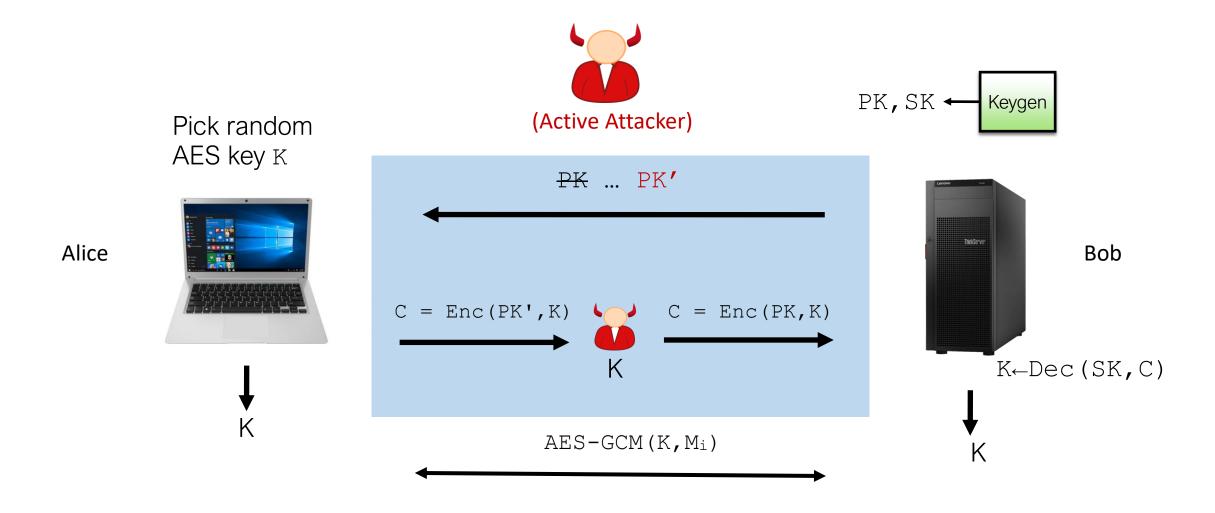
**Symmetric Encryption ("Record Protocol")** 



- Last lecture: Naïve strategy can be secure against passive adversaries.
- But the above template does <u>not</u> provide authentication & integrity.

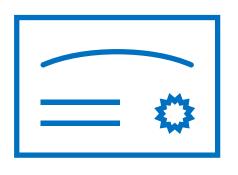
### Securing Key Exchange against Active Attackers

**Key Challenge:** Authenticity: How do we know that PK is really Bob's?



### Authentication with Certificates ("Certs")

### Suppose we had a globally trusted entity, BlaséInc.



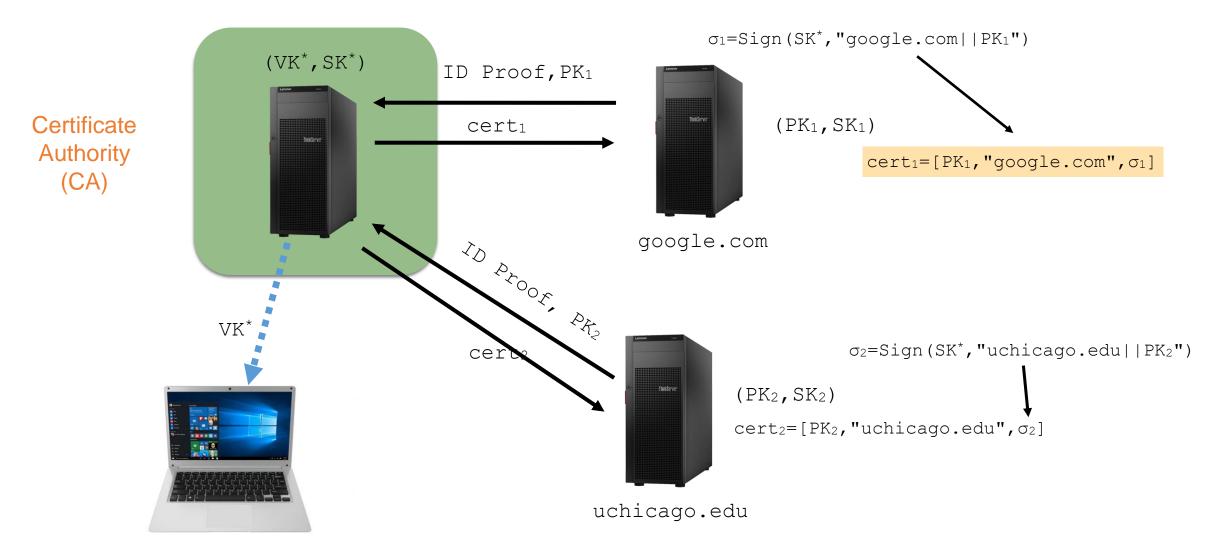
BlaséInc could issue certificates ("certs") that state what other organizations' public keys are.

Cert = a document that says:

- 1. An Entity (e.g., Bob) has a public key that is:
- 2. pk=0x7b5532..., where the document is
- 3. signed using the BlaséInc's private signing key

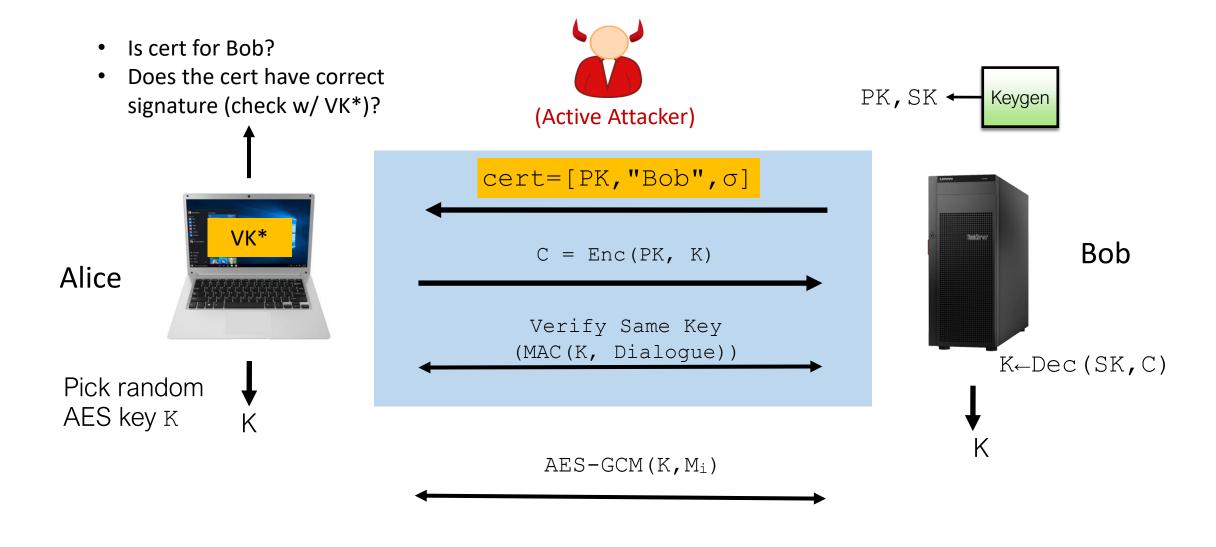
Trusted entity, BlaséInc, known as a Certificate Authority (CA)

### Authentication with Certificates ("Certs")



VK\* pre-installed on every machine by manufacturer or built into OS code.

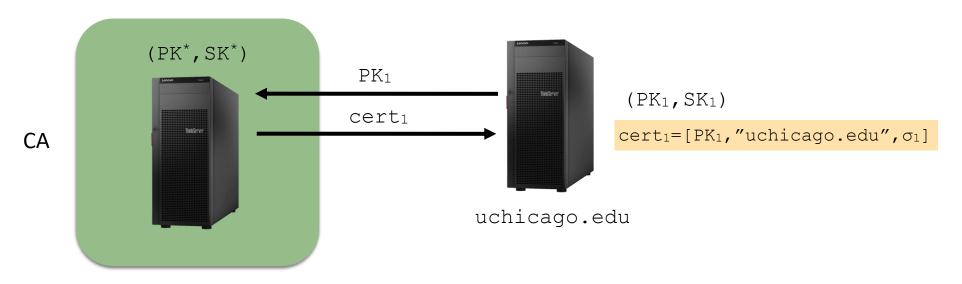
### Securing Key Exchange against Active Attackers



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### Issuing Certificates: Validation

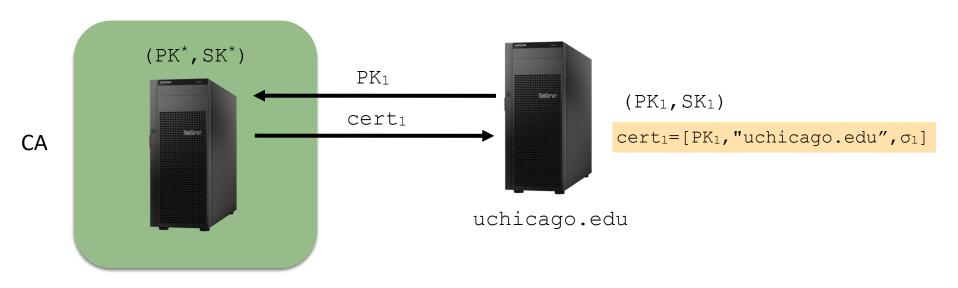


• CA must check that key PK<sub>1</sub> really does belong to "uchicago.edu"

**Domain Validation (DV)**: Check that party with that key can control domain.

**Org. Validation (OV)** and **Extended Validation (EV)**: Also check company name, location etc via public records.

### ACME Protocol by Let's Encrypt

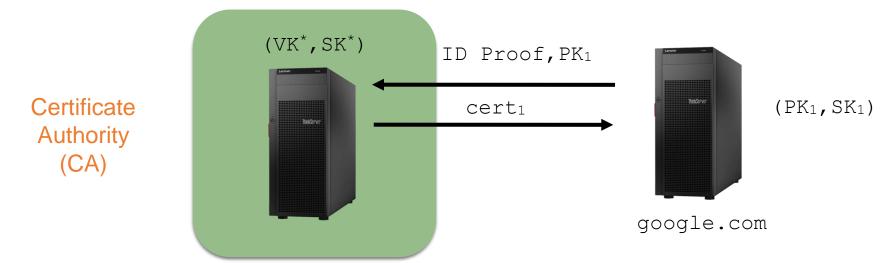


- 1. Requestor submits public key and request to CA
- 2. CA gives a *challenge* to requestor
- 3. Requestor places *challenge* on server or DNS records
- 4. CA checks *challenge* and then issues cert if challenge matches



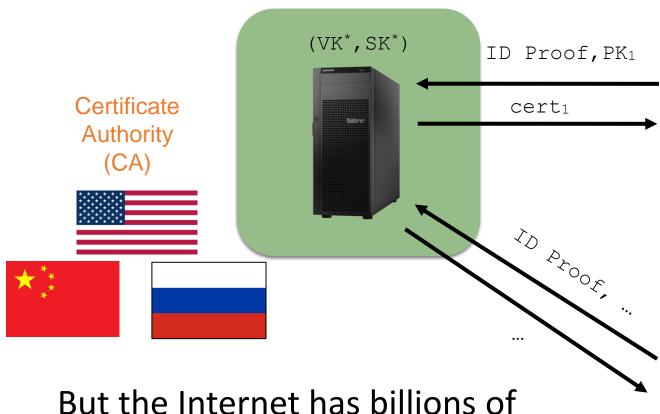


## Scaling Certificates to the Internet



Having one CA works fine if the Internet has just a few entities and everyone agrees that the CA is trustworthy.

### Scaling Certificates to the Internet

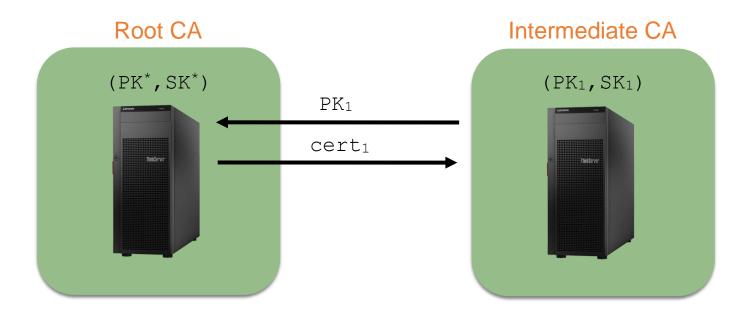


But the Internet has billions of devices...

And not everyone agrees on a trusted party (CA)...



### Scaling: Intermediate CAs and Cert Chains



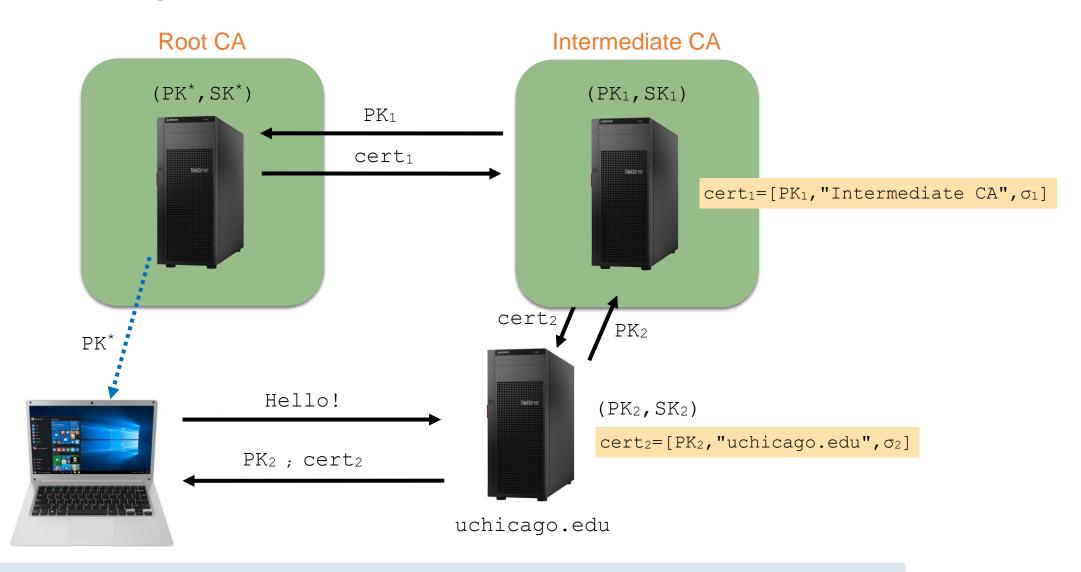
### To handle scaling:

- Allow a trusted Root CA to delegate their trust to multiple intermediate CA's
- Any of these intermediate CA's can then create a certificate for someone
  - 100's of intermediate CA's on the Internet

### Scaling: Intermediate CAs and Cert Chains

To check PK<sub>2</sub> recursive validation:

- Check cert<sub>2</sub>
  to make sure
  PK<sub>2</sub> for
  uchicago.edu
- 2) Get PK<sub>1</sub> and cert<sub>1</sub> to check sig of cert<sub>2</sub>
- 3) If cert<sub>1</sub> issued by root CA, use PK\* to its check sig.

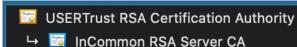


 $PK^*$  bound to Root  $\Rightarrow PK_1$  bound to CA  $\Rightarrow PK_2$  bound to uchicago.edu

### X.509 Certificates

#### Cert Content Includes:

- Cert's Serial number
- Cert's Expiration date
- Common name of subject (e.g., Bob [google.com])
- Public key of subject
- Extensions (possibly many)
- CA info (name of CA that is issuing the cert, etc.)
- CA's Signature (on hash of cert)



→ 🚟 \*.uchicago.edu

### Who are we trusting?

**Public Key Info** 

**Algorithm** RSA Encryption (1.2.840.113549.1.1.1)

**Parameters** 

None

Public Rev 256 bytes: CA E9 01 25 77 E9 74 B8 CB F7 99 DA D6 87 79 35 D7 31 CA D7 83 11 83 32 FA FA 43 CC C8 85 7B 76 EF 79 BB 4B 8B E0 35 87 EE A4 34 17 DC 5A 0D 5A 04 D3 F1 BA E7 98 9F 49 FC D5 B9 2C FB C8 DD 36 47 4D 07 FE 41 11 75 B0 42 F7 6D 40 4C BF F5 B6 C7 FE 05 0D DE 3B 7C E9 9F 6A 1C 1C 89 2E AA E8 F5 E3 5B 04 55 16 B0 48 92 C7 F9 37 11 89 F8 C5 85 C1 24 96 71 6F 78 B6 6B 35 39 92 8C EF 17 91 D1 97 D7 EF 93 6E 95 F1 EE C6 0D 5A EA 39 C6 4E 33 E2 CA F2 9A 41 F4 A2 41 9C E8 EA 46 FB EF 71 CO A6 D3 C6 A5 94 81 4B 12 5E 80 63 87 7C 2F A6 8A A5 9A 31 9E 81 63 7F 0F 26 25 B6 6D 62 C2 AD B4 E7 68 FD C9 F8 86 2C 3F F8 E1 59 F3 3E 73 08 DF 6C 92 98 21 D2 AD EF 23 E7 33 A2 D4 5E 67 74 E3 AB 08 DF 15 31 9A 9D 3B 36 7D 6B 77 48 60 17 A4 10 F3 17 77 53 E0 21 D9 F9 A4 12 0F 39 DA D1

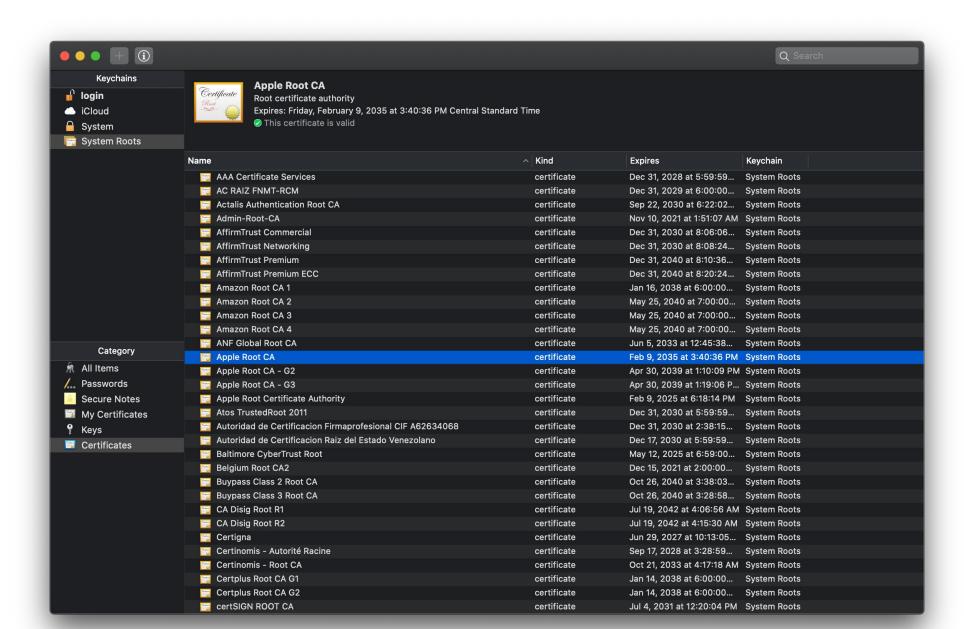
Exponent 65537

Kev Size 2.048 bits

Key Usage Encrypt, Verify, Wrap, Derive

Who's signature? Signature 256 bytes: 11 F9 F9 6D C6 92 D1 B9 E7 13 E6 0D BA E6 19 65 BB 16 4B DE E1 C2 3A 62 55 D1 61 80 93 F0 2A B2 7D 9E 76 CE 10 4A D6 96 4E 5C 00 5D BD 8C 83 74 CF C1 14 91 2B 15 4B 2D 67 4A 84 A2 A4 54 7A B1 C9 8E F5 A7 93 8D 30 BF 0C 9B EF 98 36 D6 4B BD B6 11 63 C2 51 23 71 7B 8D 4C 9B B7 AD A9 FE A8 4E 48 B2 83 A1 36 75 97 2B 36 4A 72 C4 AA C6 B6 A8 4A C0 F4 37 BD 0E 85 B1 A8 FB EC B6 B5 BB A8 C2 C0 BB B7 47 D7 D4 DB 05 80 72 BA CB C7 79 81 63 CC 55 D7 68 9C 41 2B E7 D9 F0 C2 8F 11 15 7D C5 D5 34 27 5C 7C B5 D9 A8 3F 3C DF C5 1D AA 52 03 19 AE 5B FC FF 42 68 15 A3 01 CB F8 0E FE 9B A1 76 B8 43 1C 6B 9C 57 38 87 81 3B 4A 33 98 09 CF 25 F4 75 34 AE 1E 7B CD 0F EF A0 4C 5B 92 B7 F1 FD 66 1B 49 67 B0 65 5A 90 1D 1D 54 D2 CF FF FD 07 DC 7A 88 56 51 55 16 7F 83 D4 FC 19 F4 28

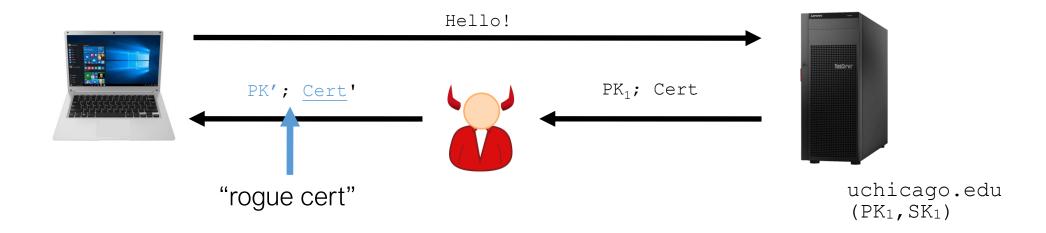
### Root CA's & Root Certificates



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## What if attacker got a "valid" cert for uchicago.edu that has their malicious key?



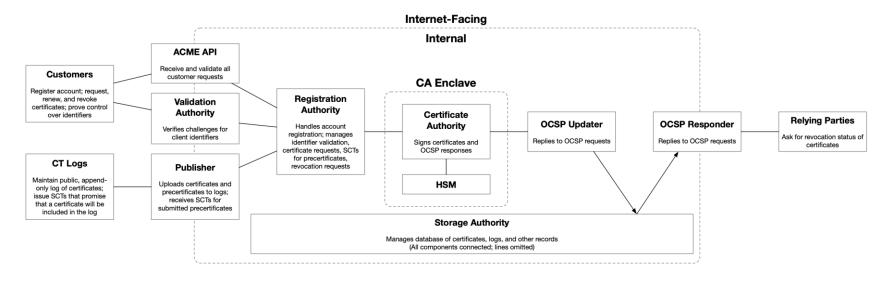
• "Machine-in-the-middle" can read/change all traffic undetected



### **CA Security**

### Some common attacks to get rogue certificate:

- Fool or bypass a CA's validation process
- Compromise a CA organization and generate malicious cert's



**Figure 3: Boulder architecture.** Let's Encrypt developed and operates a Go-based open-source CA software platform named Boulder, which is composed of single-purpose components that communicate over gRPC, as illustrated here. The certificate lifecycle unfolds roughly from left to right in the diagram.

"Let's Encrypt: An Automated Certificate Authority to Encrypt the Entire Web", CCS 2019

### Sample of CA Security Incidents

- 2011, Root CA Comodo: Login credentials stolen. Hacker issues certs for mail.google.com, login.live.com, www.google.com, login.yahoo.com...
- 2011, Root CA DigiNotar: Hacker issues rogue cert for \*.google.com, others. Used to MitM by Iranian government.
- 2013, Root CA TurkTrust: Accidentally issues intermediate CA cert, used to issue gmail.com cert.

•

• 2019, Root CA Comodo: Pushes email login credentials to public GitHub repo...

### Countermeasure: Public-Key Pinning

- Goal: Eliminate Root / Intermediate CA's with bad hygiene or who you don't trust
- Server (e.g., website) can tell client (e.g., browser) to only accept certs signed by certain CA's
  - Code trusted CA keys into client app (e.g., Chrome only trusts certs signed by Google's CA), or
  - Send special application message telling client what to pin (More common)
- Helped discover some rogue certs from previous slide
- What are some problems with this defense?
  - If server hacked... attacker can pin a malicious key/cert: will only connect w/ attacker cert!
  - Website error: pin wrong or broken key... website inaccessible!

Now deprecated because of these issues

### Countermeasure: Revocation

Publicly list bad (revoked) certificates so they are no longer accepted

• CA or Server (that was issued cert) can revoke

## Mass Revocation: Millions of certificates revoked by Apple, Google & GoDaddy

The DarkMatter debate is already having industry-wide ramifications

Millions of SSL/TLS certificates – among other digital certificates – are being revoked right now as a result of an operational error that caused the generation of non-compliant serial numbers.

## Let's Encrypt to Revoke 3 Million SSL Certificates on March 4

The world's leading free SSL provider announces that millions of certificates are being revoked due to a bug they discovered days ago — giving subscribers potentially only hours to respond

### Cert Revocation Lists (CRLs)

#### CA's CRL Server



Revoked serial numbers:

09823342365

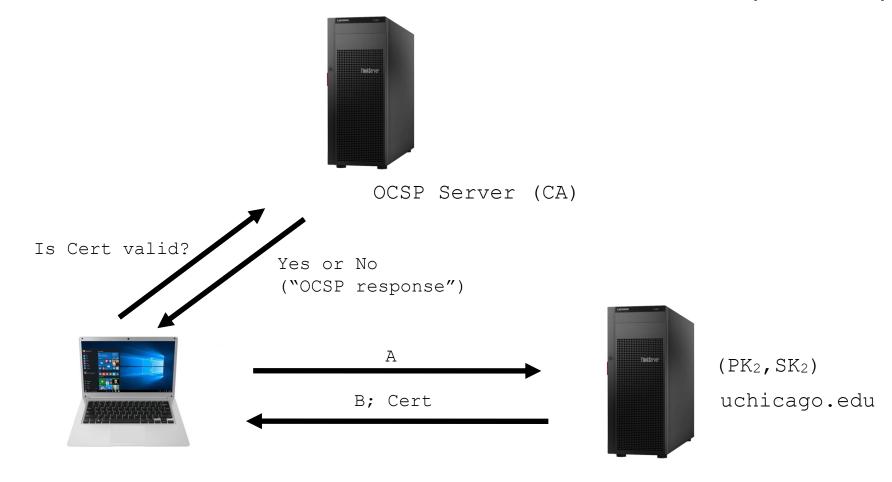
23423482349

98072344456

- Each CA provides a list of revoked cert's
- Clients can download CRL and check cert's they receive against the list
- Problems:
  - List will get too large
  - Difficult to keep current

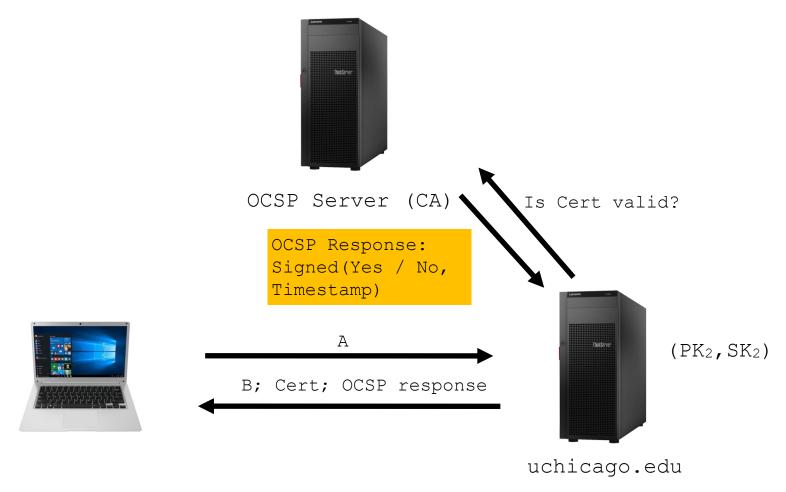


### Revocation: Online Certificate Status Protocol (OCSP)



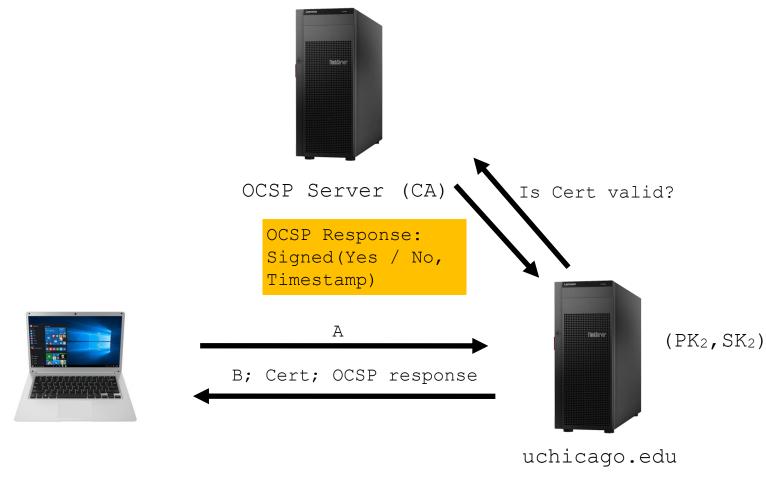
- Add another server to connect to, slowing connection
- What if OCSP server times out?
- Privacy problem?

### Revocation: OCSP Stapling



- TLS Extension that allows for OCSP response to be included with cert
- Client checks CA signature and time-stamp on response (~hours old).
- Certs can have "must staple" extension.

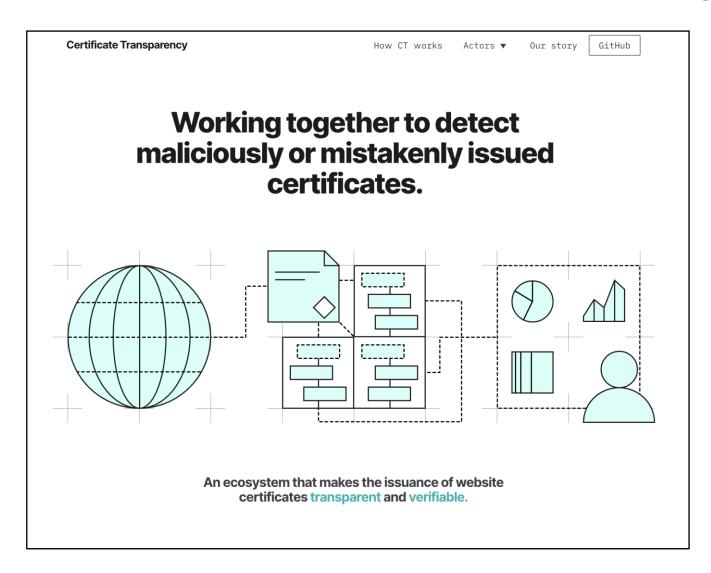
### Revocation: OCSP Stapling



Problems?

OCSP server goes down => uchicago.edu goes down (no OCSP response to attach to cert)

## Certificate Transparency (CT): How do we find rogue certs?



Scenario: Attackers compromise a CA and create rogue certs for google.com that have

- (1) attacker's public keys and
- (2) valid CA signature

How does Google or the CA discover these rogue certs were issued or in use?

#### **Cert Transparency:**

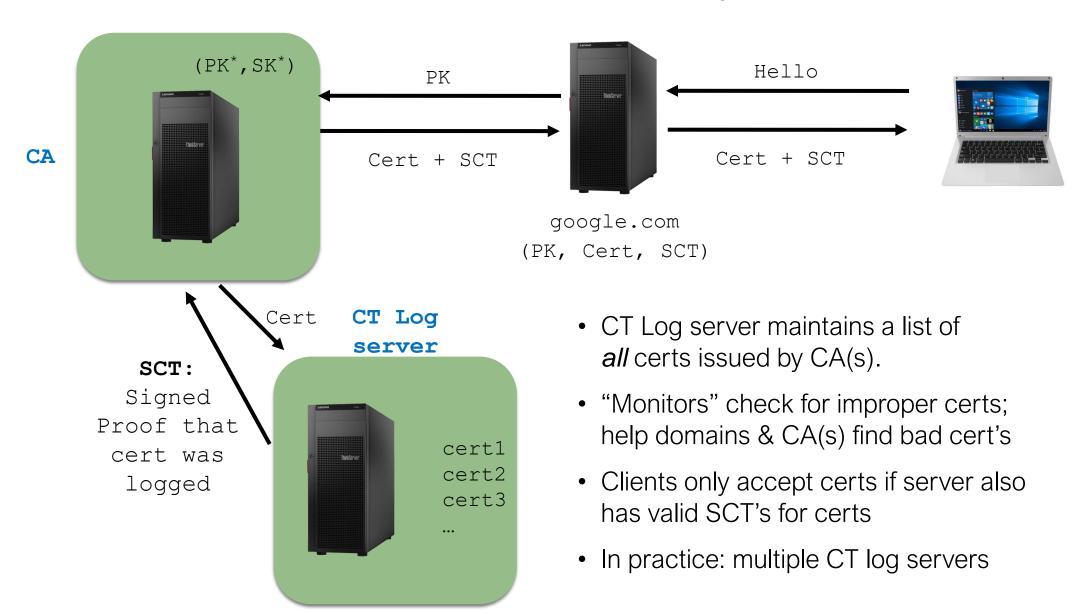
- Require all cert's added to public audit logs
- Domains & CA's can check audit logs for rogue certs & revoke them

## Certificate Transparency (CT)

### Simplified strategy to find certificates we should revoke:

- An auditor maintains a list (log) of every certificate ever issued
- Whenever a CA issues a cert, they submit (add) cert to this log
- Clients only accept a server's cert if it appears on the log
- Each server (domain) can now monitor the logs to see if anyone (and who) issued a rogue certificate for them
  - If so, add the rogue cert to revocation lists
  - If CA has pattern of issuing rogue cert's, ban them

## Certificate Transparency (CT)

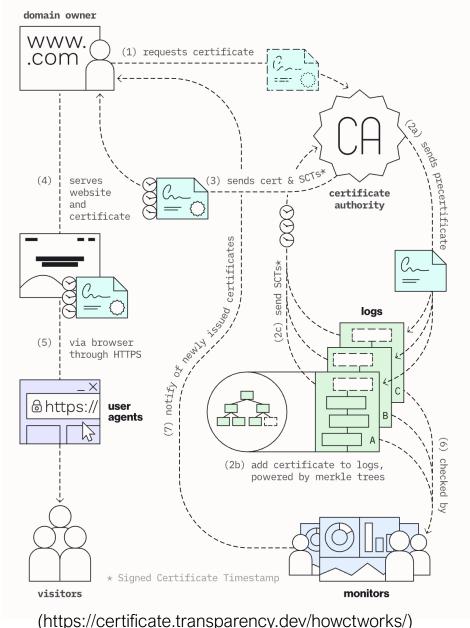


## Challenges with CT

- List is huuuuge (every issued cert... solution: temporal sharding)
- Trust the CT Log?
- (Monitors) Who checks the logs?
- Privacy (e.g., enterprise has private servers)?



## Cert Transparency & OCSP



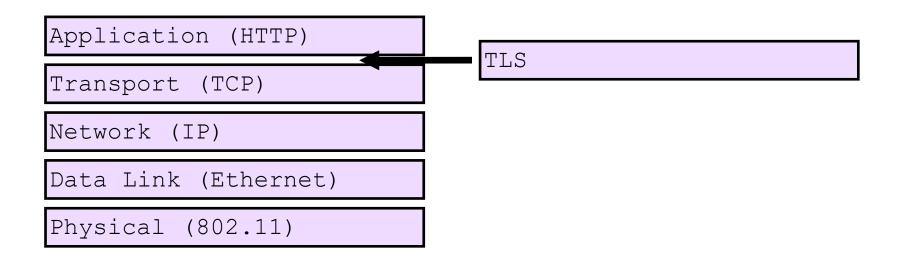
How do CT and OCSP compare?

- OCSP: Allows clients to determine if a cert is valid
- CT: Allows domains (cert owners) and CA's to find malicious cert's

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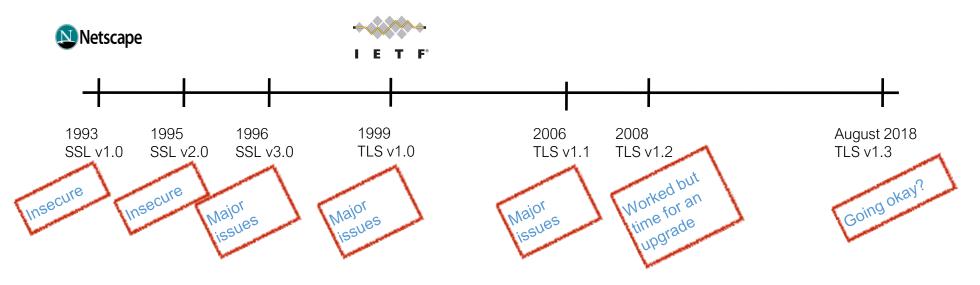
### TLS in the Protocol Stack



- Goal: Allow any application using TCP to transmit data with E2E security
- TLS takes requests from applications (e.g. browser speaking HTTP) and transmits them securely to another host on the Internet

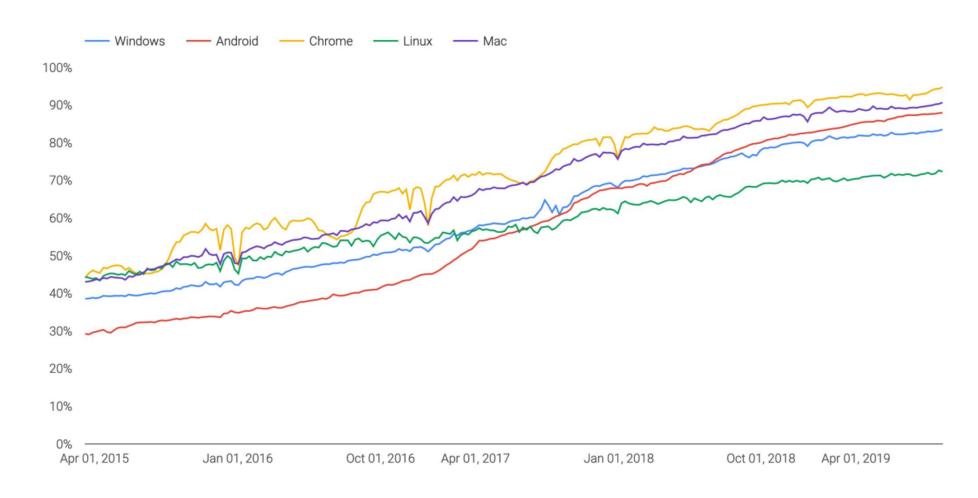
### History: SSL/TLS

- SSL = "Secure Sockets Layer"
- TLS = "Transport Layer Security" (renaming of SSL)



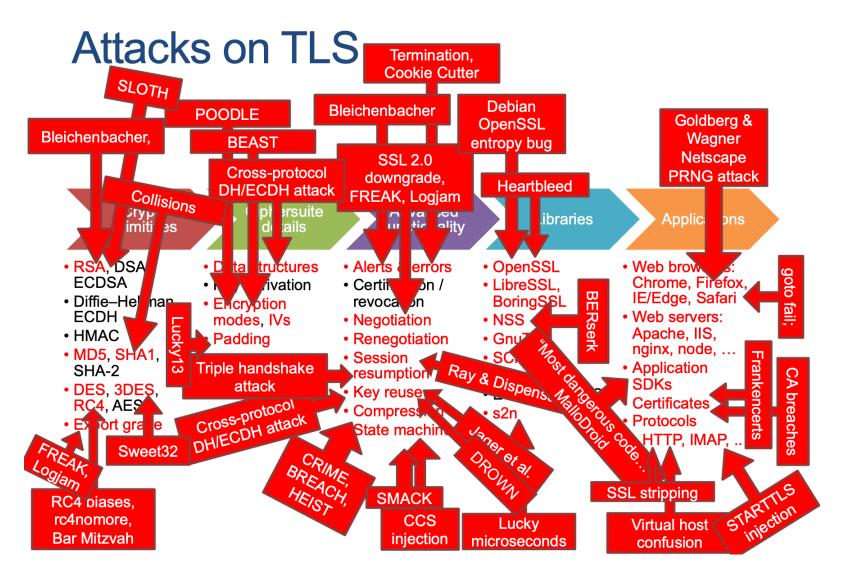
### TLS Adoption (HTTPS)

Percentage of pages loaded over HTTPS in Chrome by platform



(Source: <u>transparencyreport.google.com</u>, via Matt Green)

Attacks on TLS Stebila • 2018-09-04 5



#### Security

### It's official: TLS 1.3 approved as standard while spies weep

Now all you lot have to actually implement it

By Kieren McCarthy in San Francisco 13 Aug 2018 at 22:19 26 ☐ SHARE ▼



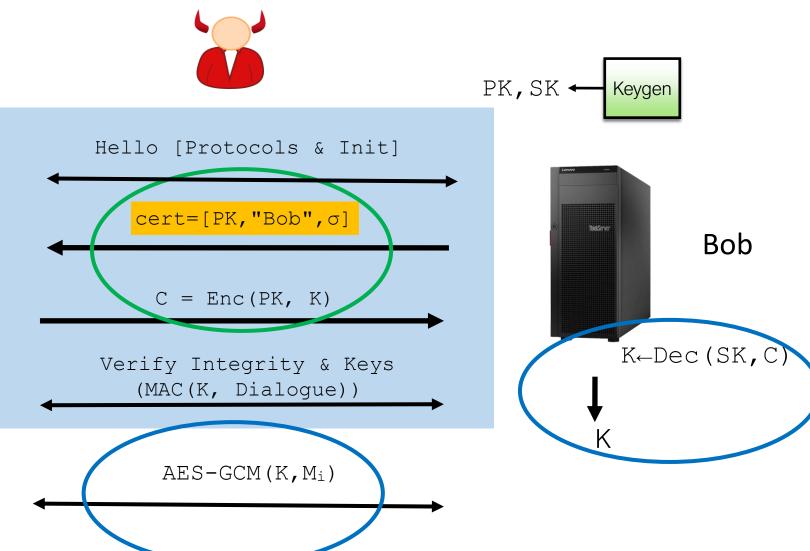
An overhaul of a critical internet security protocol has been completed, with TLS 1.3 becoming an official standard late last week.

### TLS Protocol: Very Similar to Our Template

- Is cert for Bob?
- Is cert in CT logs and has it been revoked?

 Does the certificate chain have valid signatures?





### Registrar has set Final Exam Schedule

Wed, Mar 6 from 8-10 PM

(For **BOTH** sections)

## The End