

1. Recall the discussion of polymorphic typechecking in Handout 5. Assume that we have both `int` and `real` as base types. To extend the typechecker to support overloaded functions on integers and reals (e.g., “+”), we need to allow type variables that are restricted to be members of some set. For example, the type of “+” could be written as

$$\forall \alpha \in \{\text{int}, \text{real}\}. (\alpha \times \alpha) \rightarrow \alpha$$

We can model this by changing the representation of type-variable kinds:

```
and tvar_kind
  = INSTANCE of ty
  | UNIV of int
  | NUMKIND
```

Give a modified version of the unification algorithm from Section 4 of Handout 5 that deals with this new kind representation.

2. Consider the following lexically scoped language of integer expressions:

$$exp ::= NUM \tag{1}$$

$$| VAR \tag{2}$$

$$| exp_1 \textbf{ where } VAR = exp_2 \tag{3}$$

$$| exp_1 + exp_2 \tag{4}$$

Give an attribute grammar that computes the value of an expression. You may assume that `NUM.value` is the integer value of the numeric literal and that `VAR.name` is the name of a variable. Your solution may use functional data structures, such as sets and finite maps.