

1. Consider the following lexically scoped language of integer expressions:

$exp ::= NUM$ (1)

| VAR (2)

| $\text{let } VAR = exp_1 \text{ in } exp_2$ (3)

| $exp_1 + exp_2$ (4)

Give an attribute grammar that computes the value of an expression. You may assume that $NUM.value$ is the integer value of the numeric literal and that $VAR.name$ is the name of a variable. Your solution may use functional data structures, such as sets and finite maps.

2. Consider the following representation of terms in SML:

```
datatype term = T of (string * term list)
```

where the `string` is the operator name. It is possible to define strategy combinators for this term representation, where a strategy has the type

```
type strategy = term -> term option
```

and `NONE` denotes failure. For example,

```
fun <+ (s1, s2) t = (case s1 t  
  of NONE => s2 t  
  | someT => someT  
  (* end case *))
```

implements deterministic choice and

```
fun all s (T(f, args)) = let  
  fun try ([], l) = SOME(T(f, List.rev l))  
  | try (t::ts, l) = (case s t  
    of NONE => NONE  
    | SOME t' => try(ts, t'::l)  
    (* end case *))  
  in  
    try (args, [])  
  end
```

implements the `all` combinator.

- (a) Give the SML code for the `test` combinator. Recall that the `test` combinator acts as the identity when its argument succeeds and fails when its argument fails.
- (b) Give the SML code for a generic congruence operator with the following specification:

```
val congruence : (string * strategy list) -> strategy
```