Discrete Mathematics – CMSC-37110-1 Homework 11 – November 15, 2005 Instructor: László Babai Ry-164 e-mail: laci[at]cs[dot]uchicago[dot]edu

HOMEWORK. Please print your name on each sheet. Please try to make your solutions readable. This homework is due next class, THURSDAY, NOVEMBER 17.

Please write solutions to challenge problems on a separate sheet.

READ: the ENTIRE "Graphs and Digraphs" handout (GD), especially the section on Planarity (pp 55-59) and Digraph Terminology (pp60-65).

The DO problems are for practice, do not hand them in.

- DO11.1 Prove:  $F_n = \left\lceil \left( \frac{1+\sqrt{5}}{2} \right)^n \right\rceil$ . Here  $\lceil x \rfloor$  denotes the integer nearest to x. ( $F_n$  is the n-th Fibonacci number.)
- DO11.2 Prove: if G is a planar graph with no triangles and at least 3 vertices then  $m \leq 2n 4$ . (Here n is the number of vertices and m is the number of edges.) Hint: in class we proved  $m \leq 3n 6$  for all planar graphs with at least 3 vertices. Modify the proof to take into account that there are no triangles.)
- DO11.3 Prove: a plane tree has one region.
- DO11.4 Study the Platonic solids as planar graphs (tetrahedron, cube, octahedron, dodecahedron, icosahedron). Verify Euler's formula for each.
- DO11.5 What is the difference between a plane graph and a planar graph?
- DO11.6 Prove: if a plane graph has n vertices, m edges, r regions, and c connected components then n m + r = c + 1.
- DO11.7 Draw a planar graph G such that every vertex has degree  $\geq 2$  but the graph is not Hamiltonian (has no Hamilton cycle).
- DO11.8 Prove: every bipartite graph with  $\leq 5$  vertices is planar.
- DO11.9 (a) Prove: if in a directed graph there is a directed walk from x to y then there is a directed path from x to y. (b) Prove that "accessibility" is a transitive relation on the vertices of a graph.
- DO11.10 Prove: Mutual accessibility is an equivalence relation on the set of vertices of a digraph.
- DO11.11 Prove: Accessibility defines a partial order on the set of strong components of a digraph.
- DO11.12 Prove: a digraph can be topologically sorted if and only if it is a DAG (directed acyclic graph). (Hint: prove the following lemma: a DAG always has a vertex of out-degree zero.)

Homework (due at the beginning of class Thursday, Nov 17).

- HW11.1 (5+8 points) Consider the sequence defined by the recurrence  $b_0 = 7$ ,  $b_1 = 16$ ,  $b_n = 5b_{n-1} 6b_{n-2}$ . (a) Find a simple closed-form expression for the ordinary generating function  $f(x) = \sum_{n=0}^{\infty} b_n x^n$ .
  - (b) Use this expression to find a simple explicit formula for  $b_n$  (as a linear combination of two geometric progressions).

- HW11.2 (6 points) Prove: every planar graph is 6-colorable. (Note: 6-colorability does not mean all 6 colors must actually be used. It only says 6 colors suffice.) (Hint. Use the result that every planar graph has a vertex of degree  $\leq 5$ .)
- HW11.3 (9 points) Prove: almost all graphs are not planar. (Hint: prove that it is exponentially unlikely that a given vertex has degree  $\leq 5$  in a random graph.)
- HW11.4 (10 points) GD6.2.15 (if every vertex has the same in-degree as out-degree and the digraph is weakly connected then it is strongly connected)
- HW11.5 (4 points) Prove: A DAG with n vertices has at most  $\binom{n}{2}$  edges.
- HW11.6 (10 points) Let G be a strongly connected digraph. Let a be the g.c.d. of the lengths of all cycles in G. Let x be a vertex and let b(x) be the g.c.d. of the lengths of all closed walks in G, starting at x. Prove:  $(\forall x \in V)(b(x) = a)$ .