CMSC 10500-1: Homework 5

(due on Friday July 9th)

Merge Sort

In the class we saw a way to sort a list of numbers. The basic idea there was to remove one element of the list, sort the rest of the list and insert this element into the resulting list. For obvious reasons this method is called in **Insertion Sort**. In this exercise we will see another method to sort a list of numbers, called **Merge Sort**.

This is a typical case of a **Divide and Conquer algorithm**, and consists of the following phases:

- Base Case: The list has at most 1 element.
- Divide: Split the list (approximately) evenly into two lists.
- Conquer: Recursively sort the two lists.
- Combine: Merge the two sorted lists into one sorted list.
- 1. (6 pts) To implement the divide phase, write scheme functions first-k and rest-k which consume a list and a number k and produce a list. first-k produces a list containing the first k items of the list, and rest-k contains all but the first k items of the list. If the list has at most k items then rest-k returns an empty list.

```
> (first-k (cons 1 (cons 3 (cons 2 empty))) 2)
(cons 1 (cons 3 empty))
> (rest-k (cons 1 (cons 3 (cons 2 empty))) 2)
(cons 2 empty)
```

2. (10 pts) To implement the combine phase, write a scheme function merge which consumes two lists of numbers, which are independently sorted (in increasing order) and produces one sorted list containing all the numbers in both the lists.

3. (4 pts) Finally implement the function merge-sort which consumes a list of numbers alon and sorts them. Be sure to split your list (approximately) evenly into two lists. You may use the in-built scheme function length which returns the length of a list.

Solution

Writing first-k is easy: we just need to observe that the first k elements of a list is equal to the first element and the first k-1 elements of (rest list). This together with the right boundary cases (a.k.a base cases) takes care of this. Similarly all but the first k elements of a list is equal to all but the first k-1 elements of (rest list).

```
;; first-k: list number -> list
     return list of the first k elements
(define (first-k lst k)
  (cond
    [(empty? lst) empty]
    [(= k 0) empty]
    [else (cons (first lst) (first-k (rest lst) (- k 1)))]
  )
)
;; rest-k: list number -> list
     return all but the first k elements
(define (rest-k lst k)
  (cond
    [(empty? lst) empty]
    [(= k 0) lst]
    [else (rest-k (rest lst) (- k 1))]
  )
)
```

This is just a generalization of the insert-number we did in class (if the first list has only one number, then it is the same as insert-number). Here we compare the first elements of the both the lists, and the smaller of the two becomes the first element of the output list. Then we recursively call ourselves with one of the lists getting shorter (whichever one had the smaller element).

```
;; merge: list-of-numbers list-of-numbers -> list-of-numbers
;; consume two sorted lists and output one sorted list
```

```
(define (merge lon1 lon2)
  (cond
     [(empty? lon1) lon2]
     [(empty? lon2) lon1]
     [(< (first lon1) (first lon2))
          (cons (first lon1) (merge (rest lon1) lon2))]
     [else (cons (first lon2) (merge lon1 (rest lon2)))]
)</pre>
```

Finally to sort a list of number, we solve the simple cases (less than 2 elements) directly. For other cases, we split the list into two parts using first-k and rest-k using an appropriate value of k (half the length of the list), sort the two parts recursively and merge the resulting sorted lists.

Again note that instead of choosing k to be half the length of the list, if we chose k equal to 1 always, we get insertion sort. In fact no matter what value of k we choose the algorithm sorts the list. Different values only affect the efficiency of the resulting algorithm. Choosing k as half the length of the list, results in the most efficient algorithm.